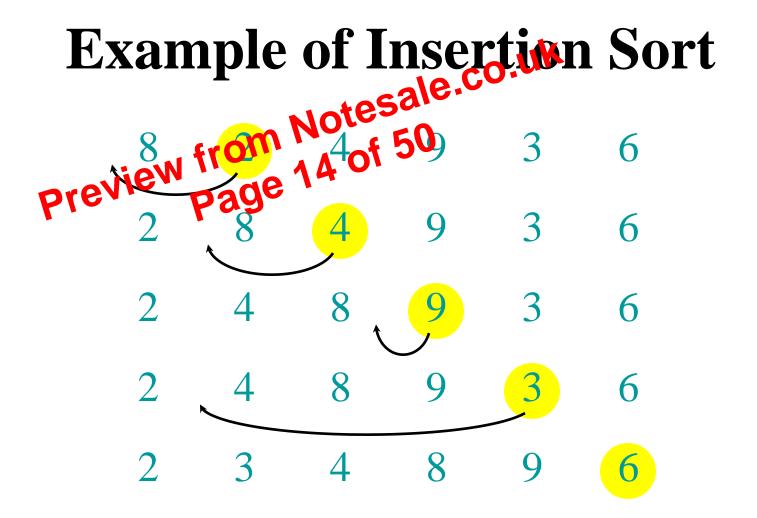
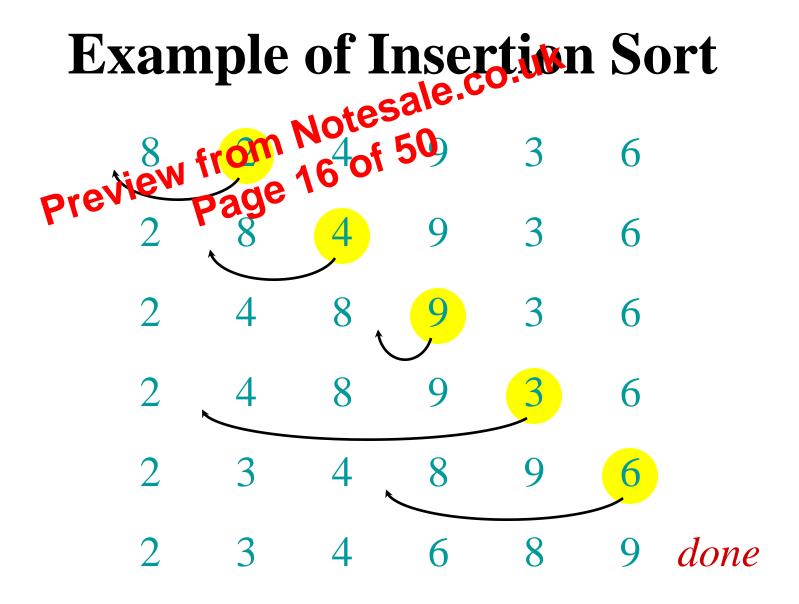
# The problem of sources Notesale. Sources Notesale. $Input Sequence (a_1, a_2, ..., a_n)$ of numbers.

**Output:** permutation  $\langle a'_1, a'_2, ..., a'_n \rangle$  such that  $a'_1 \leq a'_2 \leq \cdots \leq a'_n$ .

## Example: *Input:* 8 2 4 9 3 6 *Output:* 2 3 4 6 8 9



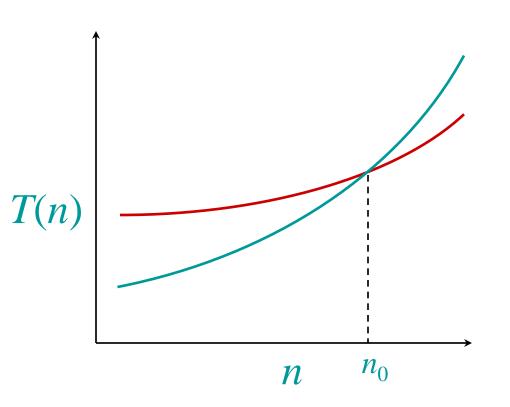


# $$\begin{split} & \bigoplus \text{-notation} \\ & Math: eview from Notesale.co.uk} \\ & Math: eview from 20 of 50 \\ & \Theta(g(n)) = \{ f(n) : \text{there exist positive constants } c_1, c_2, \text{ and} \\ & n_0 \text{ such that } 0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \\ & \text{ for all } n \geq n_0 \} \end{split}$$

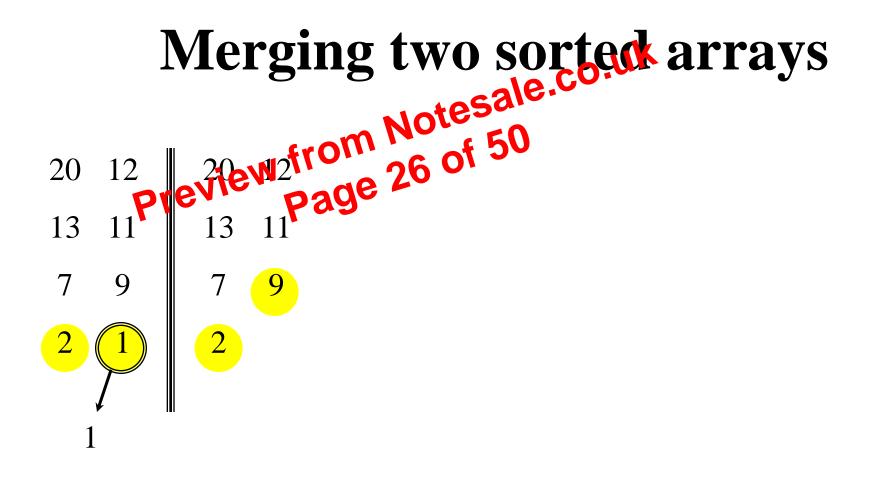
### **Engineering:**

- Drop low-order terms; ignore leading constants.
- Example:  $3n^3 + 90n^2 5n + 6046 = \Theta(n^3)$

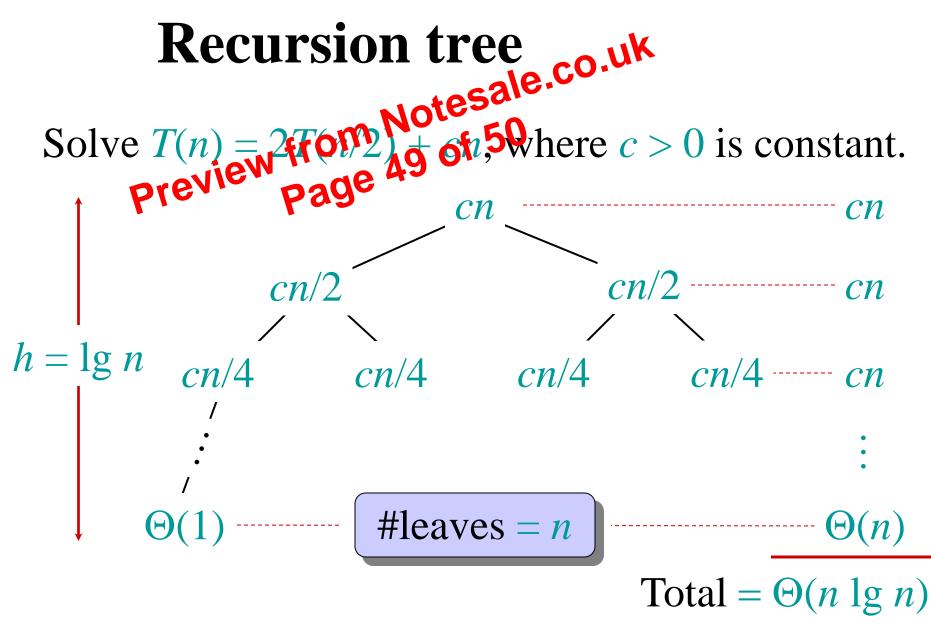
# Asymptotic performance When *n* gets large enough a $\Theta(n^2)$ algorithm *alwaye* beats a $\Theta(n^2)$ algorithm.



- We shouldn't ignore asymptotically slower algorithms, however.
- Real-world design situations often call for a careful balancing of engineering objectives.
- Asymptotic analysis is a useful tool to help to structure our thinking.



# Recursion tree Solve T(n) = 2 from bote sale couk preview age 39 of 5 where c > 0 is constant.



L1.49